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Short communication

A note on the paper 'Single machine scheduling problems with financial resource constraints: Some complexity results and properties' by E.R. Gafarov et al.

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ABSTRACT

This note emends an incorrectness in the NP-hardness proof of problem 1|NR, $d_j = d$, $g_j = g$ | $\sum T_j$ given in a paper by Gafarov et al. in Mathematical Social Sciences (see vol. 62, 2011, 7–13).

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In the article E.R. Gafarov, A.A. Lazarev, F. Werner, Single machine scheduling problems with financial resource constraints: Some complexity results and properties, Mathematical Social Sciences, 62 (2011), 7–13, the following mistake is found in Section 3.2, where the authors consider the problem denoted as 1|NR, $d_j = d$, $g_j = g$ | $\sum T_j$ and claim that it is NP-hard. In the proof, a reduction from the Partition Problem was used which is not polynomial, since M exponentially depends on n .

However, it is not difficult to correct this proof. The main idea of using M^{n-i+1} was that the processing time of a job

belongs to a pair with the smallest number being greater than the total sum of the processing times of all jobs from the pairs with larger numbers, e.g., for the job V_2 : $p_2 \gg \sum_{i=2}^n (p_{2i-1} + p_{2i})$.

Instead of using $p_{2i} = M^{n-i+1}$, where $M = (n \sum b_j)^n$ (see the definition of the instance given in (3) on page 11), we can consider, e.g., $p_{2i} = 2n \cdot 2^{n-i+1} M$, where $M = (n \sum b_j)$. In this case, the reduction will be polynomial in the input length, if we suppose that all digits used are coded in a binary system with approximately 2^n zero–one symbols per digit.

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